







Advantages w/ Remote Exploder

- The mailing list can stay anonymous to the sender
- Maybe good for bandwidth (imagine all members of a mailing list is in Mars)
- Save bandwidth if the mailing list is very long
- Can be in parallel when multiple mailing lists

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Advantages w/ Local Exploder

- Easier to prevent mail forwarding loops
- Duplicate copy prevention
- Can estimate bandwidth cost before sending out emails

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Establishing Keys
Establishing public keys

Out of band mechanism
PKI
Piggybacking certificates on emails

Establishing secret keys

Alice phones Bob . . . (we knows this is bad)
Kerberos

Privacy

- Why?
 - Eavesdropper
 - Relay nodes (routers or MTAs)
- · End-to-end privacy
- · Privacy with distribution list exploders

End-to-end Privacy

- · Alice sends Bob an email that is encrypted with Bob's public key
- Well, not ideal, because
 - Multiple recipients
 - Public key crypto is far less efficient than secret key crypto
 - Better not to use long term key unless really the only way to do so

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A Public Key Based E2E Privacy Solution

• Alice picks up a secret key and then sends out the following:

> Bob's name; $K_{Bob}{S}$ Carols's name; $K_{carol}{S}$ Ted's name; $K_{Ted}{S}$ $S\{m\}$

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Authentication of the Source

- · Source authentication based on public key technology
- Sign the message using the sender's private key · Source authentication based on secret keys
 - A message must carry a MAC (message authentication code)
 - MAC can be:
 - · CBC residue of the message computed with the shared secret key
 - · Message digest of the shared secret append to the message • Encrypted message digest (preferred when multiple recipients)
- · Source authentication with distribution lists

Message Integrity

- Source authentication often must come with the message integrity
 - Otherwise, why care the source authentication?
- But how about message integrity w/o source authentication?
 - Can be done if the message is encrypted with the recipient's public key
 - Perhaps needed by a kidnapper

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Non-Repudiation

- · Non-repudiation based on public key technology - Relatively easy
 - Require the message to be signed by the sender using its private key
 - · Remember nobody else knows the private key, so . . .
- · Non-repudiation with secret keys
 - Relatively difficult
 - The message is signed using a shared secret key But nobody else knows the secret key (what's the difference here from above?)

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Plausible Deniability Based on Public Key Technology

• Alice picks a secret key S

- $\{S\}_{\text{Bob}}$ (encrypted with Bob's public key)
- [{S}_{Bob}]_{Alice} (signed with Alice's private key)
- *MAC* of m = f(S, m)
- Alice sends the following to Bob:
 - $m, MAC, [\{S\}_{Bob}]_{Alice}$
- Bob can know that *m* is from Alice, but he can't prove to anyone else that m is from Alice Copyright Bobs can know S



Proof of Submission
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Anonymity	

	Containment
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Verifying WHEN a message was really sent

- Preventing Backdating
- · Preventing Postdating

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Quiz 4

- Assume secret key crypto. If Bob wants to verify that an email is indeed from Alice, he will check a piece of data that comes with the message:
- (1) What's that piece of data called?
- (2) Who calculated this piece of data?
- (3) List three different ways to calculate this piece of data.

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Main Goals of PEM

- Privacy
- Integrity
- Source authentication
- PEM uses the similar methods we talked earlier

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PEM Model

- Smart PEM software sitting at the source and the destination
- User keys are used to sign or encrypt - One key per message
- User keys are based on either secret key or public key technology

PEM Message Structure

- A PEM message can contain several parts
- · And each part treated differently
 - Clear text
 - Integrity protected
 - Or encrypted
- With markers around each block

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Types of Message Pieces

- Ordinary, unsecured data
- MIC-CLEAR
- Clear text + MIC
- MIC-ONLY
- Encoded text + MICENCRYPTED
 - Encoded (Encrypted (clear text) + encrypted(MIC))
- Note: MIC here is the PEM's term for MAC

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PEM Certificate Hierarchy

- · A hierarchy of CAs in a tree form
 - The root is called **IPRA** (Internet Policy Registration Authority)
 - CAs certified by IPRA are called PCA (Policy Certificate Authority)
 - Then other CAs
- Policy: each CA has a policy on issuing certificates
 - Three different policies

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CA Types

- High Assurance (HA) CA
 - Super secureVery strict on deciding to issue a certificate to
 - somebody
- Discretionary Assurance (DA) CA
 Well managed, but no guarantee
- No Assurance (NA) CA
 No constraints as long as no duplications

















Private Key

- Needed when
 - Signing your own message
 - Decrypting a message delivered to you that is encrypted using your public key
- PGP can generate a private key for you – Then store it in an encrypted form

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Midterm

- 7-10 problems
- 2-3 essay questions
- 80 minutes
 - 2 3:20 p.m. Nov 18th
- First couple weeks are covered in course reserve materials
- Use the lecture slides as the guidance - Textbook and course reserves as reference
- The level of materials details to remember - To the level that slides have
 - But not to the level of textbooks
- Know steps of Kerberos, SSL, ...