

Another Proof of the Completeness Theorem

The first three exercises provide a proof of completeness. We have broken this proof up into three parts to aid understanding and reduce complexity.

1. Prove the following Lemma A:

Lemma A:

$$\phi_1, \phi_2, \dots, \phi_n \models \psi \quad \Rightarrow \quad \models \phi_1 \rightarrow (\phi_2 \rightarrow (\dots (\phi_n \rightarrow \psi) \dots))$$

2. Prove Lemma B:

Lemma B:

Consider a propositional formula ϕ with propositional variables p_1, p_2, \dots, p_n . Let T_ϕ be the truth table of ϕ , and let L be the line number of a row of this truth table.

For all $1 \leq i \leq n$, define $\hat{p}_i = p_i$ if the entry in line L for p_i is 1, and $\hat{p}_i = \neg p_i$ if the entry in line L for p_i is 0.

Then the following holds:

$$\hat{p}_1, \hat{p}_2, \dots, \hat{p}_n \vdash \phi \text{ if the truth value (final column) of line } L \text{ is 1.}$$

$$\hat{p}_1, \hat{p}_2, \dots, \hat{p}_n \vdash \neg\phi \text{ if the truth value (final column) of line } L \text{ is 0.}$$

3. Prove the Completeness Theorem for propositional logic:

Completeness Theorem:

$$\Gamma \models \phi \quad \Rightarrow \quad \Gamma \vdash \phi$$

First Order Logic

4. Consider the sentence

$$\phi = \forall x \exists y \exists z (P(x, y) \wedge P(z, y) \wedge P(x, z) \rightarrow P(z, x)).$$

Which of the following models satisfies ϕ ?

- a The model \mathcal{M} consists of the set of natural numbers with $P^{\mathcal{M}} = \{(m, n) \mid m < n\}$.
- b The model \mathcal{M}' consists of the set of natural numbers with $P^{\mathcal{M}'} = \{(m, 2*m) \mid m \text{ is a natural number}\}$.
- c The model \mathcal{M}'' consists of the set of natural numbers with $P^{\mathcal{M}''} = \{(m, n) \mid m < (n + 1)\}$.

5. Let ϕ be the formula

$$\forall x \forall y \exists z (R(x, y) \rightarrow R(y, z))$$

where R is a predicate symbol of two arguments.

- (a) let $A = \{a, b, c\}$ and $R^{\mathcal{M}} = \{(b, c), (b, b), (b, a)\}$. Do we have $\mathcal{M} \models \phi$?
- (b) let $A = \{a, b, c\}$ and $R^{\mathcal{M}'} = \{(b, c), (a, b), (c, b)\}$. Do we have $\mathcal{M}' \models \phi$?