



# A Simple Pacemaker Implementation



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• A medical device that regulates the beating of the heart using electrical impulses to contract the heart muscles.

• The primary purpose of the pacemaker is to maintain an adequate heart rate.

 Modern pacemakers are externally programmable and allow the cardiologist to select the optimum pacing modes for individual patients.





 Implement behavior described in the Pacemaker Requirements Document by Boston Scientific.

- Need to demonstrate:
  - Utility : it performs helpful actions
  - Safety : it does not perform harmful actions.



• The pacemaker function modes are described by simple mnemonics.

	I	II		IV(optional)
Category	Chambers Paced	Chambers Sensed	Response To Rate	Rate Modulation
Letters	O–None A–Atrium V–Ventricle D–Dual	O–None A–Atrium V–Ventricle D–Dual	O–None T–Triggered I–Inhibited D–Tracked	R–Rate

#### Modes implemented:

- AAT: a very simple mode that "paces" the atrium every time a "sense" is detected.
- VVI: a mode that paces the ventricle depending on whether there is or is not a natural ventricle sense.
- DDD: a more complex mode that paces both the atrium and/or ventricle depending on atrial and ventricular senses.



 Formal specification of a subset of the natural language requirements.

- TRIO used as the specification language.
- Analysis of the formal specification to check consistency.
  - PVS and ZOT used.
- Simulation of the pacemaker behavior.
  - Implemented in Java on a desktop platform.
- Simulation tested.

Solution



#### • The project is an example of:

- A safety critical system
- A real-time system



- Use of temporal first order logic to specify the behavior of the system enables us to produce test cases very easily and simplify the design.
- We were also able to prove the correctness of some properties: safety and utility.





TRIO

• TRIO is a Formal Language based on a metric extension of first-order temporal logic and exploits typical object-oriented features to support the managing of large, complex, and maintainable specifications.

 The basic operator Dist(F,t) specifies that F holds after/before t time instants.



page 8



#### **TRIO Operators**

- Futr(F, d)  $\Leftrightarrow$  d  $\geq$  0  $\land$  Dist(F, d)
  - F will be true after a time interval of d.
- Lasts[ee|ii|ie|ei](F, d) ⇔ ∀d' (0<d'<d→Dist(F, d'))</li>
  *F* holds over a period of length d.
- Lasted[ee|ii|ie|ei](F, d)  $\Leftrightarrow \forall d'(0 < d' < d \rightarrow Dist(F, -d'))$

F held over a period of length d in the past.

• NowOn(F)  $\Leftrightarrow \exists d (d>0 \land Lasts(F,d))$ 

F holds over a certain period of unspecified length.





Boston Scientific Requirements:

- •The system consists of 3 components
  - Pulse Generator (PG)
  - Device Control Monitor (DCM)
  - Leads

## • TRIO+ Modular Specification:

• Using the concept of classes it is possible to group together sets of axioms that refer to the same component.





VVI

### V: Ventricle Paced

V: Ventricle Sensed

 I: "A sense in a chamber shall inhibit a pending pace in that chamber"

## • TRIO Specification:

- AXIOM Alw( sensesignalV AND NOT ignoresignalV IFF senseV )
- 2) AXIOM Alw( senseV IMPLIES NowOn(ignoresignalV) )
- 3) AXIOM Alw(VVI IMPLIES

(Lasted(NOT senseV,TIMEOUT) IFF artpulseV))



#### AXIOM Alw( sensesignalV AND NOT ignoresignalV IFF senseV )

senseV is an event occurring when there is a signal in the ventricle and it is not ignored.

2) AXIOM Alw( senseV IMPLIES NowOn(ignoresignalV) )

Detected the ventricular event senseV then ignore any signal for a certain time interval (refractory period).





VVI







VVI





## **ZOT: Sat Checker**

- The ZOT tool is used to verify the satisfiability of the overall system TRIO axiomatization.
- If a contradiction is found, a counterexample is shown.
- Problem: The check is performed at discrete time.



Not enough for the proof of crucial system properties



time 0 VVI	time 9 ARTPULSEV VVI	time 16 VVI
time 1 VVI	IGNORESIGNALA_ARP SENSESIGNALA PWAVE	time 17 VVI
time 2		time 18
VVI	time 10	VVI
	VVI	NATPULSEA
time 3	SENSEV	
VVI	IGNORESIGNALA_ARP	time 19
	SENSESIGNALV	**LOOP**
time 4	IMPONV	VVI
VVI		SENSEA
	time 11	SENSESIGNALA
time 5	VVI	PWAVE
VVI	IGNORESIGNALV	
	IGNORESIGNALA_PVARP	time 20
time 6	SENSESIGNALV	ARTPULSEV
VVI	IMPONV	VVI
		IGNORESIGNALA_ARP
time 7	time 12	SENSESIGNALA
VVI	VVI	PWAVE
NATPULSEA	IGNORESIGNALV	
	IGNORESIGNALA_PVARP	time 21
time 8		VVI
VVI	time 13	SENSEV
SENSEA	VVI	IGNORESIGNALA_ARP
SENSESIGNALA		SENSESIGNALV
PWAVE	time 14	IMPONV
	VVI	
	time 15	

VVI

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A Simple Pacemaker Implementation



 It is a Theorem Prover: a tool that automates some typical logical operation necessary for a formal demonstration.

- Human support required.
- Supports the TRIO Axiomatization. TVS: TRIO PVS
- When a property is proved, it holds for all the possible models even in continuous time.



The property is valid



• **Utility:** If the patient Heart Rate (HR=1/RR) is not within the normal range the pacemaker has to pace the heart artificially as defined by the required LRL. Note that Timeout= 1/LRL.

CONJECTURE

Alw( senseV AND RR>Timeout AND VVI AND Lastsii(VVI, Timeout) IMPLIES

Futr(senseV, Timeout) AND Lasts(NOT senseV, Timeout))

• **Safety:** If the heart behaves normally the pacemaker does not interfere with its natural pulse. CONJECTURE

Alw( senseV AND RR<=Timeout AND VVI AND Lastsii(VVI, RR) IMPLIES

Futr(senseV, RR) AND Lasts(NOT senseV, RR))



PVS

- ZOT
  - Automatically checks the satisfiability of the system.
  - Produces examples useful in understanding the correctness of the axiomatization.
- PVS
  - Used to prove the target properties.
  - Forces a manual process that clarifies which axiom is not correct and why not.

Every time a new axiom is introduced, this validation process is repeated.



Simulation is needed to:

- Have a visual verification of the Pacemaker behavior, validating the simulated results with the expected one in the analysis steps.
- Develop a visual understanding of the different function modes.
- Demonstrate how the formal approach we used is well suited to rapidly produce a software prototype.





#### **Java Simulator**

 Java classes represent all the Pacemaker modules:

- Device Control Monitor, to set Pacemaker parameters.
- Heart Behavior (not presented in the BS Requirements).
- Pulse Generator set in the appropriate functioning mode.
- Electrocardiogram as the output of the simulation.

 TRIO axioms are manually translated into Java methods.





#### **Java Simulator**

#### • Heart parameters:

- HAV(atrial-ventricular interval).
- RR (natural heart period).
- Wave Amplitudes.
- Wave Durations.

#### Pacemaker Parameters

- Modes.
- System Timeouts.
- Refractory periods.

La sut Deve a star				$\mathbb{Z}$
Heart Parameters		Pacen	naker Parameters	
HAV		-		
300	•	Funct	ioning Mode	
RR		VVI		-
1000 -				
PWAVEDUR		AV TI	meout	
100	Ŧ			
PWAVEAMP		200		
2 🗸		LRL Timeout		
QRSWAVEDUR				_
250	-	1200		-
QRSWAVEAMP			afractory Bariad	
6 🗸		VRP Refractory Period		
ARTWAVEDUR_A		300 -		
100 👻				
ARTWAVEAMP_A		ARP F	lefractory Period	
2	-	300		-
ARTWAVEDUR_V		1		
250	•	PVAR	P Refractory Period	1
ARTWAVEAMP_V		300		-
6	•			
		C	ommands	
Ouick Sample	He	art	Start Simulation	



 At the beginning the heart has a natural pace which is regular and admissible.

 After a while the device is called to perform ventricular artificial pulses to resume the heart rate (HR) over the lower rate limit (LRL).





## Conclusions

• The use of the TRIO language allowed us to produce a formal specification:

- Readable and understandable.
- Compact.
- Supported by powerful tools.
- The approach proposed by the complementary use of ZOT and PVS guarantees:
  - A full understanding of the requirements.
  - A very fast implementation step.
- The simulation:
  - Reflects the behavior expected in the analysis step thanks to very usable graphs.