

Statistical Debugging

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Bug Isolation Architecture



Winnowing Down the Culprits

- 1710 counters
 - 3 × 570 call sites
- 1569 zero on all runs
 141 remain
- 139 nonzero on at least one successful run
- Not much left!
 - file_exists() > 0
 - > xreadline() == 0





Multiple, Non-Deterministic Bugs

- Strict process of elimination won't work
 - Can't assume program will crash when it should
 - No single common characteristic of all failures
- Look for general correlation, not perfect prediction
- Warning! Statistics ahead!



Consider each predicate P one at a time

▶ Include inferred predicates (e.g. \leq , \neq , \geq)

- ▶ How likely is failure when *P* is true?
 - (technically, when P is observed to be true)
- Multiple bugs yield multiple bad predicates



F(P) = # failing runs with |P| > 0 S(P) = # successful runs with |P| > 0 $Bad(P) = \frac{F(P)}{S(P) + F(P)}$



Are We Done? Not Exactly!

$$Bad(f = NULL) = 1.0$$



Are We Done? Not Exactly!

$$Bad(f = NULL) = 1.0$$
$$Bad(x = 0) = 1.0$$

Predicate (x = 0) is innocent bystander Program is already doomed



Identify unlucky sites on the doomed path

$$Context(P) = \frac{F(P \lor \neg P)}{S(P \lor \neg P) + F(P \lor \neg P)}$$

 Background risk of failure for reaching this site, regardless of predicate truth/falsehood



Does *P* being true *increase* the chance of failure over the background rate?

Increase(P) = Bad(P) - Context(P)

Formal correspondence to likelihood ratio testing



Increase Isolates the Predictor

Increase(f = NULL) = 1.0

Increase(x = 0) = 0.0



... for programs with just one bug.

- Need to deal with multiple bugs
 - How many? Nobody knows!
- Redundant predictors remain a major problem

Goal: isolate a single "best" predictor for each bug, with no prior knowledge of the number of bugs.



Multiple Bugs: Some Issues

- A bug may have many redundant predictors
 - Only need one, provided it is a good one
- Bugs occur on vastly different scales
 - Predictors for common bugs may dominate, hiding predictors of less common problems



- Multiple interesting & useful predicate metrics
- Simple visualization may help reveal trends



Confidence Interval on Increase(P)

$$\pm 1.96 \sqrt{\frac{Bad(P) \cdot (1 - Bad(P))}{S(P) + F(P)}} + \frac{Context(P) \cdot (1 - Context(P))}{S(P \lor \neg P) + F(P \lor \neg P)}$$

Strictly speaking, this is slightly bogus

- Bad(P) and Context(P) are not independent
- Correct confidence interval would be larger



Bad Idea #1: Rank by *Increase(P)*

Thermometer	Context	Increase	S	F	F + S	Predicate
	0.065	0.935 ± 0.019	0	23	23	((*(fi + i)))->this.last_token < filesbase
	0.065	0.935 ± 0.020	0	10	10	((*(fi + i)))->other.last_line == last
	0.071	0.929 ± 0.020	0	18	18	((*(fi + i)))->other.last_line == filesbase
	0.073	0.927 ± 0.020	0	10	10	((*(fi + i)))->other.last_line == yy_n_chars
	0.071	0.929 ± 0.028	0	19	19	bytes <= filesbase
	0.075	0.925 ± 0.022	0	14	14	((*(fi + i)))->other.first_line == 2
	0.076	0.924 ± 0.022	0	12	12	((*(fi + i)))->this.first_line < nid
	0.077	0.923 ± 0.023	0	10	10	((*(fi + i)))->other.last_line == yy_init

- High Increase but very few failing runs
- These are all sub-bug predictors
 - Each covers one special case of a larger bug
- Redundancy is clearly a problem



Bad Idea #2: Rank by *F(P)*

Thermometer	Context	Increase	S	F	F + S	Predicate	
	0.176 0.176 0.176 0.176	$\begin{array}{c} 0.007 \pm 0.012 \\ 0.007 \pm 0.012 \\ 0.007 \pm 0.012 \\ 0.007 \pm 0.013 \end{array}$	22554 22566 22571 18894	5045 5045 5045 4251	27599 27611 27616 23145	<pre>files[filesindex].language != 15 tmp == 0 is FALSE strcmp != 0 tmp == 0 is FALSE</pre>	
	0.176	0.007 ± 0.013	18885	4240	23125	files[filesindex].language != 14	
	0.176	0.008 ± 0.013	17757	4007	21764	filesindex >= 25	
	0.177	0.008 ± 0.014	16453	3731	20184	new value of M < old value of M	
	0.176	0.261 ± 0.023	4800	3716	8516	<pre>config.winnowing_window_size != argc</pre>	

- Many failing runs but low Increase
- Tend to be super-bug predictors
 - Each covers several bugs, plus lots of junk



A Helpful Analogy

In the language of information retrieval

- Increase(P) has high precision, low recall
- F(P) has high recall, low precision

Standard solution:

- Take the harmonic mean of both
- Rewards high scores in both dimensions



Rank by Harmonic Mean

Thermometer Context	Increase	S	F	F + S	Predicate	
0.176	0.824 ± 0.009	0	1585	1585	files[filesindex].language > 16	
0.176	0.824 ± 0.009 0.824 ± 0.009	0	1584 1580	1584 1580	strcmp > 0 strcmp == 0	
0.176	0.824 ± 0.009	0	1577	1577	files[filesindex].language == 17	
0.176	0.824 ± 0.009	0	1576	1576	tmp == 0 is TRUE	
0.176	0.824 ± 0.009 0.883 ± 0.012	1	774	775	$((*(fi + i))) \rightarrow this.last line == 1$	
0.116	0.883 ± 0.012	1	776	777	((*(fi + i)))->other.last_line == yyleng	

Definite improvement

- Large increase, many failures, few or no successes
- But redundancy is still a problem



Redundancy Elimination

One predictor for a bug is interesting

- Additional predictors are a distraction
- Want to explain each failure once
- Similar to minimum set-cover problem
 - Cover all failed runs with subset of predicates
 - Greedy selection using harmonic ranking



- 1. Rank all predicates under consideration
- 2. Select the top-ranked predicate P
- 3. Add *P* to bug predictor list
- 4. Discard *P* and all runs where *P* was true
 - Simulates fixing the bug predicted by P
 - Reduces rank of similar predicates
- 5. Repeat until out of failures or predicates



Case Study: Moss

Reintroduce nine historic Moss bugs

- High- and low-level errors
- Includes wrong-output bugs
- Instrument with everything we've got
 - Branches, returns, variable value pairs, the works
- 32,000 randomized runs at ¹/₁₀₀ sampling



Effectiveness of Filtering

Scheme	Total	Retained	Rate
branches	4170	18	0.4%
returns	2964	11	0.4%
value-pairs	195,864	2682	1.4%



Effectiveness of Ranking

- Five bugs: captured by branches, returns
 - Short lists, easy to scan
 - Can stop early if Bad drops down
- Two bugs: captured by value-pairs
 - Much redundancy
- Two bugs: never cause a failure
 - No failure, no problem
- One surprise bug, revealed by returns!



Analysis of exif



- 3 bug predictors from 156,476 initial predicates
- Each predicate identifies a distinct crashing bug
- All bugs found quickly using analysis results



Analysis of Rhythmbox

Initial	Effective	Predicate
		<pre>tmp is FALSE (mp->priv)->timer is FALSE (view->priv)->change_sig_queued is TRUE (hist->priv)->db is TRUE rb_playlist_manager_signals[0] > 269 (db->priv)->thread_reaper_id >= 12 entry == entry fn == fn klass > klass genre < artist vol <= (float)0 is TRUE (player->priv)->handling_error is TRUE (statusbar->priv)->library_busy is TRUE shell < shell len < 270</pre>

- 15 bug predictors from 857,384 initial predicates
- Found and fixed several crashing bugs



How Many Runs Are Needed?

		Fai	ling Ru	uns For	⁻ Bug #	n	
	#1	#2	#3	#4	#5	#6	#9
Moss	18	10	32	12	21	11	20
ccrypt	26						
bc	40						
Rhythmbox	22	35					
exif	28	12	13				



Other Models, Briefly Considered

- Regularized logistic regression
 - S-shaped curve fitting
- Bipartite graphs trained with iterative voting
 - Predicates vote for runs
 - Runs assign credibility to predicates
- Predicates as random distribution pairs
 - Find predicates whose distribution parameters differ
- Random forests, decision trees, support vector machines, ...



Capturing Bugs and Usage Patterns

- Borrow from natural language processing
 - Identify topics, given term-document matrix
 - Identify bugs, given feature-run matrix
- Latent semantic analysis and related models
 - ► Topics ⇔ bugs and usage patterns
 - Noise words ⇔ common utility code
 - Salient keywords ⇔ buggy code



Probabilistic Latent Semantic Analysis





Cluster runs by most probable topic
Failure diagnosis for multi-bug programs

- Characterize representative run for cluster
 - Failure-inducing execution profile
 - Likely execution path to guide developers
- Relate usage patterns to failure modes
 - Predict system (in)stability in scenarios of interest



Compound Predicates for Complex Bugs



"Logic, like whiskey, loses its beneficial effect when taken in too large quantities."

> Edward John Moreton Drax Plunkett, Lord Dunsany, "Weeds & Moss", *My Ireland*



Limitations of Simple Predicates

Each predicate partitions runs into 2 sets:

- Runs where it was true
- Runs where it was false

Can accurately predict bugs that match this partition

- Unfortunately, some bugs are more complex
 - Complex border between good & bad
 - Requires richer language of predicates



Motivation: Bad Pointer Errors





Motivation: Bad Pointer Errors

Kinds of Predicate	Best Predicate	Score
Simple Only	new len == old len	0.71
Simple & Compound	o + s > buf_size ∧ offset < len	0.94

Too many compound predicates

- 2^{2^N} functions of *N* simple predicates
- N² conjunctions & disjunctions of two variables
- N ~ 100 even for small applications
- Incomplete information due to sampling
- Predicates at different locations



Conservative Definition

- A conjunction $C = p_1 \land p_2$ is true in a run iff:
 - p_1 is true at least once and
 - p_2 is true at least once
- Disjunction is defined similarly
- Disadvantage:
 - C may be true even if p_1 , p_2 never true simultaneously
- Advantages:
 - Monitoring phase does not change
 - ▶ $p_1 \land p_2$ is just another predicate, inferred offline



Three-Valued Truth Tables

For each predicate & run, three possibilities:

- 1. True (at least once)
- 2. Not true (and false at least once)
- 3. Never observed

Conjunction: $p_1 \wedge p_2$

Disjunction: $p_1 \lor p_2$





Compute score of each conjunction & disjunction

- $C = p_1 \wedge p_2$
- $D = p_1 \vee p_2$

Compare to scores of constituent simple predicates

- Keep if higher score: better partition between good & bad
- Discard if lower score: needless complexity
- Integrates easily into iterative ranking & elimination



Still Too Many

• Complexity: $N^2 \cdot R$

- N = number of simple predicates
- R = number of runs being analyzed
- 20 minutes for N ~ 500, R ~ 5,000
- Idea: early pruning optimization
- Compute upper bound of score and discard if too low
 "Too low" = lower than constituent simple predicates
- Reduce O(R) to O(1) per complex predicate



Upper Bound On Score

► ↑ Harmonic mean

$$\uparrow Increase(C) = \frac{\uparrow F(C)}{\uparrow F(C) + \downarrow S(C)} - \frac{\downarrow F(C \text{ obs})}{\downarrow F(C \text{ obs}) + \uparrow S(C \text{ obs})}$$

$$\uparrow Sensitivity(C) = \frac{\uparrow \log F(C)}{\log NumF}$$

- Upper Bound on $C = p_1 \land p_2$
 - Find $\uparrow F(C)$, $\downarrow S(C)$, $\downarrow F(C \text{ obs})$ and $\uparrow S(C \text{ obs})$
 - > In terms of corresponding counts for p_1 , p_2



$\uparrow F(C)$ and $\downarrow S(C)$ for conjunction





▶ \downarrow S(C): true runs are disjoint



either 0 or

$$S(p_1)+S(p_2) - NumS$$

(whichever is maximum)



↑S(C obs)

- Maximize two cases
- C = true
 - True runs of p₁, p₂ overlap
- C = false
 - False runs of p_1 , p_2 are disjoint





Usability

- Complex predicates can confuse programmer
 - Non-obvious relationship between constituents
 - Prefer to work with easy-to-relate predicates
- *effort*(p₁, p₂) = proximity of p₁ and p₂ in PDG
 PDG = CDG ∪ DDG
 - Per Cleve and Zeller [ICSE '05]
- Fraction of entire program
 - "Usable" only if *effort* < 5%</p>
 - Somewhat arbitrary cutoff; works well in practice



Evaluation: What kind of predicate has the top score?





Evaluation: Effectiveness of Pruning



Analysis time: from ~20 mins down to ~1 min

Prune: effort > 5%

- Prune: score upper bound too low
- Compute exact score, but too low
- Compute exact score and retain

Evaluation: Impact of Sampling





Evaluation: Usefulness Under Sparse Sampling



