Scalable Defect Detection

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Part III
Lightweight Specifications for Win32 APIs

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The Win32 API

Win32 API is the layer on which all modern Windows applications are built. .NET is built on top, and contains many managed classes that wrap Win32 functionality.
Business Goals

I. Significantly reduce the number of exploitable buffer overruns in Windows Vista

II. Change development process so products after Vista are more secure
Standard Annotation Language

• Created in summer June 2002 joint effort with product groups and CSE
• Specifies programmer intent which leads to:
  – Better coverage (reduce false negatives)
  – Reduced noise (reduce false positives)
  – Ecosystem of tools
  – High impact results
Measured Outcomes

<table>
<thead>
<tr>
<th></th>
<th>Mutable String Arguments</th>
<th>Functions</th>
</tr>
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<tr>
<td>Total</td>
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</tr>
<tr>
<td>Annotated</td>
<td>1031</td>
<td>6,918</td>
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</table>

Headers for toy application only expose $\frac{1}{5}$th of all Win32 APIs

Developers did more than the minimum required for security!

```c
#include<tchar.h>
#include<windows.h>
#include<wincrypt.h>
#include<wininet.h>
#include<shlwapi.h>
#include<shlobj.h>

int _tmain(...)
{
    return 0;
}
```
How We Got There!

Manual Annotation

SAL

SALInfer

MIDL Compiler

espX, IO, MSRC, Prefast, Prefix, Truscan, ...

Warnings

Drive these to zero!

Triage Warnings

SAL Fixes

Code Fixes

Lock in progress!
Focus of This Talk

- Code
- SALInfer
- MIDL Compiler
- Manual Annotation
- SAL
- espX, IO, MSRC, Prefast, Prefix, Truscan, ...
- Warnings
- Triage Warnings
- Code Fixes
- SAL Fixes
Technical Design Goals

• Improves coverage and accuracy of static tools
• Locks in progress for the future
• Usable by an average windows developer
• Cannot break existing Win32 public APIs or force changes in data-structures (i.e. no fat pointers)
Technical Design Non-Goals

• No need to guarantee safety
• No need to be efficiently checked as part of normal foreground “edit-debug-compile” loop
• No need to handle all the corner cases
• No need to be “pretty”
Take a Peek Yourself!

For MSDN documented Win32 APIs start here


or Google “Live Search” for them

Annotated headers can be download from

Vista SDK

first search hit for “Vista SDK”
Run-Time Library Reference

**memcpy, wmemcpy**

Copies bytes between buffers. These functions are deprecated because more secure versions are available; see  

```c
void *memcpy(
    void *dest,
    const void *src,
    size_t count
);

wchar_t *wmemcpy(
    wchar_t *dest,
    const wchar_t *src,
    size_t count
);
```

**Parameters**

- **dest**
  - New buffer.
- **src**
  - Buffer to copy from.
- **count**
  - Number of characters to copy.

**Return Value**

The value of `dest`.

**Remarks**

`memcpy` copies `count` bytes from `src` to `dest`; `wmemcpy` copies `count` wide characters (two bytes). If the source and destination overlap, the behavior of `memcpy` is undefined. Use `memmove` to handle overlapping regions.

**Security Note**

Make sure that the destination buffer is the same size or larger than the source buffer. For more information, see [Avoiding Buffer Overruns](http://msdn2.microsoft.com/en-us/library/ms717795(VS.80).aspx).

The `memcpy` and `wmemcpy` functions will only be deprecated if the constant `_CRT_SECURE_DEPRECATE_MEMORY` is defined prior to the inclusion statement in order for the functions to be deprecated, such as in the example below:
memcpy, wmempcpy, (cont)

<table>
<thead>
<tr>
<th>Parameters</th>
<th>dest</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>New buffer.</td>
</tr>
</tbody>
</table>

| src | Buffer to copy from. |

| count | Number of characters to copy. |

- **Return Value**
  - The value of dest.

- **Remarks**
  - **memcpy** copies count bytes from src to dest; **wmempcpy** copies count wide characters (two bytes). If the source and destination overlap, the behavior of **memcpy** is undefined. Use **memmove** to handle overlapping regions.

- **Security Note**
  - Make sure that the destination buffer is the same size or larger than the source buffer. For more information, see [Avoiding Buffer Overruns](http://msdn2.microsoft.com/en-us/library/ms717795(VS.80).aspx).

The **memcpy** and **wmempcpy** functions will only be deprecated if the constant **_CRT_SECURE_DEPRECATE_MEMORY** is defined prior to the inclusion statement, in order for the functions to be deprecated, such as in the example below:

Not all the information is relevant to buffer overruns.

Just say “No” to bad APIs.

For every API there’s usually a wide version. Many errors are confusing “byte” versus “element” counts.
International Features

**MultiByteToWideChar**

Maps a character string to a wide character (Unicode UTF-16) string. The character string mapped by this function is not necessarily from a multibyte character set.

```c
int MultiByteToWideChar(
    UINT CodePage,
    DWORD dwFlags,
    LPCSTR lpMultiByteStr,
    int cbMultiByte,
    LPWSTR lpWideCharStr,
    int cchWideChar
);
```

This unfortunately is a typical Win32 API
int MultiByteToWideChar(
    UINT CodePage,
    DWORD dwFlags,
    LPCSTR lpMultiByteStr,
    int cbMultiByte,
    LPWSTR lpWideCharStr,
    int cchWideChar
);

\textit{lpMultiByteStr}
\begin{itemize}
  \item [in] Pointer to the character string to convert.
\end{itemize}

\textit{cbMultiByte}
\begin{itemize}
  \item [in] Size, in bytes, of the string indicated by the \textit{lpMultiByteStr} parameter. Alternatively, this parameter can be set to -1 if the string is null-terminated. Note that, if \textit{cbMultiByte} is 0, the function fails.
  \item If this parameter is -1, the function processes the entire input string, including the null terminator. Therefore, the resulting wide character string has a null terminator, and the length returned by the function includes the terminating null character.
  \item If this parameter is set to a positive integer, the function processes exactly the specified number of bytes. If the provided size does not include a null terminator, the resulting wide character string is not null-terminated, and the returned length does not include the terminating null character.
\end{itemize}

\textit{lpWideCharStr}
\begin{itemize}
  \item [out] Pointer to a buffer that receives the converted string.
\end{itemize}

\textit{cchWideChar}
\begin{itemize}
  \item [in] Size, in WCHAR values, of the buffer indicated by \textit{lpWideCharStr}. If this value is 0, the function returns the required buffer size, in WCHAR values, including any terminating null character, and makes no use of the \textit{lpWideCharStr} buffer.
\end{itemize}

\textbf{Return Values}

Returns the number of WCHAR values written to the buffer indicated by \textit{lpWideCharStr} if successful. If the function succeeds and \textit{cchWideChar} is 0, the return value is the required size for the buffer indicated by \textit{lpWideCharStr}.
How to Solve a Problem like `MultiByteToWideChar`?

• Start with an approximate specification
• See how much noise and real bugs you find
• Power up the tools and refine until you find the next thing to worry about
• Need conditional null termination to handle case when `cbMultiByte` is -1
• Buffer size weakening handles `cbMultiByte` 0 case
BCryptResolveProviders

The **BCryptResolveProviders** function obtains a collection of all of the providers that meet the specified criteria.

```c
NTSTATUS WINAPI BCryptResolveProviders(
    LPCWSTR pszContext,
    ULONG dwInterface,
    LPCWSTR pszFunction,
    LPCWSTR pszProvider,
    ULONG dwMode,
    ULONG dwFlags,
    ULONG* pcbBuffer,
    PCRYPT_PROVIDER_REFS* ppBuffer
);
```
A common pattern to communicate buffer sizes to callee

Optional reference argument!

Optional buffer!

A common pattern to communicate buffer sizes to callee
GetEnvironmentStrings

Retrieves the environment variables for the current process.

LPTCH WINAPI GetEnvironmentStrings(void);

Parameters

This function has no parameters.

Return Value

If the function succeeds, the return value is a pointer to the environment block of the current process.
If the function fails, the return value is NULL.

Remarks

The GetEnvironmentStrings function returns a pointer to a block of memory that contains the environment variables of the calling process. Each environment block contains the environment variables in the following format:

Var1=Value1\0
Var2=Value2\0
Var3=Value3\0
...
VarN=ValueN\0\0

Double null termination!
Does Your Head Hurt Yet?
Does Your Head Hurt Yet?

If only C had exceptions, garbage collection, and a better string type the Win32 APIs would be much simpler!
Does Your Head Hurt Yet?

I WISH IT DID!
The Next Best Thing

Use the .NET Win32 bindings until it does!
The Next Best Thing

So when are they going to rewrite Vista in C#?
So That’s Why It Took Five Years!

Read up about the "Longhorn Reset"
So That’s Why It Took Five Years!

Intel and AMD will solve this problem eventually! Until then we have **SAL**.
WINBASEAPI
int WINAPI MultiByteToWideChar(
    __in  UINT   CodePage,
    __in  DWORD  dwFlags,
    __in_bcount(cbMultiByte) LPCSTR  lpMultiByteStr,
    __in  int    cbMultiByte,
    __out_ecount_opt(cchWideChar) LPWSTR  lpWideCharStr,
    __in int    cchWideChar);

NTSTATUS WINAPI BCryptResolveProviders(
    __in_opt LPCWSTR pszContext,
    __in_opt ULONG dwInterface,
    __in_opt LPCWSTR pszFunction,
    __in_opt LPCWSTR pszProvider,
    __in ULONG dwMode,
    __in ULONG dwFlags,
    __inout ULONG* pcbBuffer,
    __deref_opt_inout_bcount_part_opt(*pcbBuffer, *pcbBuffer)
    PCRYPT_PROVIDER_REFS *ppBuffer);
GetEnvironmentStrings

WINBASEAPI
__out
__nullnullterminated
LPCH
WINAPI
GetEnvironmentStrings(
    VOID
);

End of Section A

Questions?
From Types to Program Logics a Recipe for SAL

A story inspired by true events
A Recipe for SAL

1) Start with a simple Cyclone like type system
2) Slowly shape it into a powerful program logic for describing common Win32 APIs
3) Add some syntactic sugar and abstraction facilities
4) Mix in a lot of developer feedback
5) Bake it until it’s properly done!

It’s getting there but still needs some cooking!
Types vs Program Logic

• Types are used to describe the representation of a value in a given program state
• Program Logic describe transitions between program states

Aside: Each execution step in a type-safe imperative languages preserves types so types by themselves are sufficient to establish a wide class of properties without the need for program logic
Concrete Values

Scalars

'\0',...'a','b','c', ...
..., -2, -1, 0, 1, 2, ...

Pointers

Cells

Extent

'\0','a', 1

'\0' 'e' 'l' 'l' 'o'
Abstract Values

A, B, ..., X, Y, Z

Some Scalar

Some Pointer

Some Cell

Some Extent

n
Program State

Roots

\[ x \rightarrow 'a' \]
\[ y \rightarrow 1 \]
\[ p \rightarrow \quad 1 \]

Store
Well-Typed Program State

char x \rightarrow \ 'a' \\
int y \rightarrow \ 1 \\
int* p \rightarrow \ 1
Well-Typed Program State

char x \rightarrow 'a'
int y \rightarrow 1
int* p \rightarrow \bullet
Well-Typed Program States

C types not descriptive enough to avoid errors
Well-Typed Program States

Use Cyclone style qualifiers to be more precise!
Generalizing `@numelts`

@numelts(3) int* buf → 

![Diagram showing three elements in memory]

int cbuf → N

@numelts(cbuf) int* buf → 

![Diagram showing N elements in memory]

What's wrong with this?
Is it Initialized or Not?

int* buf → 1 2 3

int* buf → 1 ? ?

int* buf → 1 ? 3
Define @numelts(e) as @extent(e,e)

@extent(3,3) int* buf →  

@extent(2,3) int* buf →  

@extent(??,3) int* buf →  

Just give up here!
Refined Abstract Extent

\[ \text{where } m \leq n \]
Some Special Cases

- Fully initialized extent
  - when $m = n$

- Some allocated extent
  - when $m = 0$
@extent(count, capacity)

@extent(0,3) int* buf →

int cbuf →

@extent(0, cbuf) int* buf →
Qualified Types Useful for Win32 APIs

t ::= int | void | char | t* | q₁ ... qₙ t
q ::= @range(e₁, e₁) | @relop(e,op)
    | @nonnull | @nullable | @null | @readonly
    | @numelts(e) | @allocated(e) | @extent(e₁, e₂)
    | @bsize(e) | @balloeced(e) | @bextent(e₁, e₂)
    | @zeroterm | @zerozeroterm
op ::= == | <= | >= | !=
e ::= ....
A Qualified Type for memcpy

@nonnull @numelts(count)
void* memcpy(
    @nonnull @allocated(count)
    void *dest,
    @readonly @nonnull @numelts(count)
    const void *src,
    size_t count)

It seems to work? What's wrong?
Which One is Right?

void f(@nonnull @allocated(1) int *p) {
    *p = 1;
}

void f(@nonnull @numelts(1) int *p) {
    *p = 1;
}

Types don’t capture the state transition!
Program State Transitions

Pre-condition

@allocated(1) int* p →

Post-condition

@numelts(1) int* p →

Pre-post pair make a up a contract!
Contracts with Program Logics

```c
void f( @Pre{ @nonnull @allocated(1) } 
    @Post{ @nonnull @numelts(1) } 
    int *p) {
    *p = 1;
}
```
void f(  @Pre{  @nonnull @allocated(1) }  
  @Post{  @numelts(1) }  
  int *p) {
  *p = 1;
}

Simplify because C is call by value!
Contracts with Program Logics

```c
void f( @Pre{ @nonnull @allocated(1) } @Post{ @numelts(1) }
        int *p) {
    *p = 1;
}
```

Who in their right mind is going to write that!
Contracts with Program Logics

```c
#define __out \n @Pre{ @nonnull @allocated(1) } \n @Post{ @numelts(1) }

void f(__out int *p) {
    *p = 1;
}
```

C Preprocessor macros to the rescue! Defined to empty string for compatibility.
Single Element Contracts

#define __in \ 
 @Pre{ @readonly @nonnull @numelts(1) }

#define __out \ 
 @Pre{ @nonnull @allocated(1) } \ 
 @Post{ @numelts(1) }

#define __inout \ 
 @Pre{ @nonnull @numelts(1) } \ 
 @Post{ @numelts(1) }
Single Element Contracts

#define __in_opt \
@Pre{ @readonly @nullable @numelts(1) }

#define __out_opt \
@Pre{ @nullable @alloced(1) } \
@Post{ @numelts(1) }

#define __inout_opt \
@Pre{ @nullable @numelts(1) } \
@Post{ @numelts(1) }
Contracts for Element Extents

#define __in_ecount(e) \ 
@Pre{ @readonly @nonnull @numelts(e) } 

#define __out_ecount_part(cap,count) \ 
@Pre{ @nonnull @allocated(cap) } \ 
@Post{ @extent(count,cap) } 

#define __inout_ecount_part(cap,count) \ 
@Pre{ @nonnull @extent(count,cap) } \ 
@Post{ @extent(count,cap) } 

Note order of args
# define __out_ecount_full(e) \ 
   __out_ecount_part(e,e)
# define __inout_ecount_full(e) ...
/* opt versions */
# define __in_ecount_opt(e) ...
# define __out_ecount_part_opt(cap,count) ...
# define __inout_ecount_part_opt(cap,count) ...
# define __out_ecount_full_opt(e) ...
# define __inout_ecount_full_full_opt(e) ...
Contracts for Byte Extents

#define __in_bcount(e) ...
#define __out_bcount_part(cap,count) ...
#define __inout_bcount_part(cap,count) ...
#define __out_bcount_full(e) ...
#define __inout_bcount_full(e) ...
#define __in_bcount_opt(e) ...
#define __out_bcount_part_opt(cap,count) ...
#define __inout_bcount_part_opt(cap,count) ...
#define __out_bcount_full_opt(e) ...
#define __inout_bcount_full_opt(e) ...
Developers can learn a small set of macros and be productive quickly.

<table>
<thead>
<tr>
<th>annotation</th>
<th>% total</th>
<th>% cum</th>
</tr>
</thead>
<tbody>
<tr>
<td>__in</td>
<td>47.45%</td>
<td>47.45%</td>
</tr>
<tr>
<td>__out</td>
<td>10.37%</td>
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<td>__in_opt</td>
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<td>_in_z</td>
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<td>__RPC__out_ecount_full</td>
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<td>0.21%</td>
<td>94.82%</td>
</tr>
<tr>
<td>__reserved</td>
<td>0.20%</td>
<td>95.01%</td>
</tr>
</tbody>
</table>

Distribution of macros used across Vista source base.
Contract for memcpy

__out_bcount_full(count)
void* memcpy(
  __out_bcount_full(count)
  void *dest,
  __in_bcount(count)
  const void *src,
  size_t count);

Ignore meaningless pre-condition
What about pointers to pointers?

```c
void f(__out (@nullable int*)*)* p) {
    static int l = 3;
    if(...) *p = NULL;
    else *p = &l;
}
void f(__deref_out_opt int **p) {
...
}
```

Syntax makes applying *automatically inferred* annotations to legacy code tractable!
How We Got There!

Inference bootstrapped everything!

Code

SALInfer

MIDL Compiler

SAL

Manual Annotation

espX, IO, MSRC, Prefast, Prefix, Truscan, ...

Warnings

Triage Warnings

Code Fixes

SAL Fixes
What about Nested Pointers?

#define __derefer_out_opt \
  @Pre{ @nonnull @allocated(1) } \ 
  @Deref @Post { @nullable @numelt(1) }

Pushes context of assertion down a pointer level
Annotated Types for Win32 APIs

\[ t ::= \text{int} | \text{char} | \text{void} | t^* | t \]
\[ \text{at} ::= a_1 \ldots a_n t \]
\[ p ::= \@\text{range}(e_1, e_n) | \ldots | \@\text{zerozero}\text{term} \]
\[ a ::= \@\text{Deref} a | \@\text{Pre} \{ p_1 \ldots p_n \} | \@\text{Post} \{ p_1 \ldots p_n \} | p \]
\[ \text{op} ::= \ldots \]
\[ e ::= \ldots \]

Actual primitive syntax is different. Just use the macros! Your code will be non-portable if you don't!
What About This Case?

```c
bool f(__out_opt int *p) {
    if(p != NULL) {
        *p = 1;
        return true;
    }
    return false
}
```

Need to introduce conditional contracts!
Adding `__success(cond)`

- Most conditional behavior is related to error handling protocols (i.e. exceptions via return codes)
- Introduce specialized construct for this case: `__success(expr) f(...);` means Post-conditions only hold when "expr" is true (non-zero) on return of function.
- Full conditional support on the roadmap!
Using Success

```c
__success(return == true)
bool f(__out_opt int *p) {
    if(p != NULL) {
        *p = 1;
        return true;
    }
    return false
}
```

Is _opt the right thing?
Using Success Correctly!

__success(return == true)

bool f(__out int *p) {
    if(p != NULL) {
        *p = 1;
        return true;
    }
    return false
}

Annotate for successful case!
StringCchCat Function

StringCchCat is a replacement for strcat [ http://msdn.microsoft.com/library/en-us/vclib/html/_crt_strcat.2c_.wscat.2c_.mbscat.asp ]. The size, in characters, of the destination buffer is provided to the function to ensure that StringCchCat does not write past the end of this buffer.

Syntax

```c
HRESULT StringCchCat(
    LPTSTR pszDest,
    size_t cchDest,
    LPCTSTR pszSrc
);
```

Parameters

- **pszDest**
  [in, out] Pointer to a buffer containing the string to which pszSrc is concatenated, and which contains the entire resultant string. The string at pszSrc is added to the end of the string at pszDest.

- **cchDest**
  [in] Size of the destination buffer, in characters. This value must equal the length of pszSrc plus the length of pszDest plus 1 to account for both strings and the terminating null character. The maximum number of characters allowed is STRSAFE_MAX_CCH.

- **pszSrc**
  [in] Pointer to a buffer containing the source string that is concatenated to the end of pszDest. This source string must be null-terminated.
Contracts For StringCchCat

HRESULT StringCchCat(  
  __post __nullterminated __out  
  LPTSTR pszDect,  
  __range(0,STRSAFE_MAX_CCH)  
  size_t cchDest,  
  __nullterminated __in  
  LPCTSTR pszSrc);

Much more verbose than we'd like!
Types with Contracts For StringCchCat

typedef __nullterminated TCHAR* LPSTR;
typedef const LPSTR LPCSTR;
typedef __range(0, STRSAFE_MAX_CCH) size_t STRSIZE;

HRESULT StringCchCat(
    __out LPTSTR pszDect,
    __in STRSIZE cchDest,
    __in LPCSTR pszSrc);

Must mean null terminated only in post condition!
New primitive @valid

typedef @zero termin TCHAR* LPSTR;

void f( @Pre{@nonnull @allocated(1)}
    @Post{@valid @numelts(1)}
    LPSTR s) {
    s[0]='\0';
}

Annotations associated with types only happen when an extent is "valid"
Memory Semantics Revisited

- **Allocated**
  - ?
  - Can be written to but nothing is known about its contents

- **Initialized**
  - The contents are in a known state

- **Valid**
  - Type specific properties hold
Lifecycle of a LPTSTR

Allocated

@allocated(3) LPTSTR s →

Initialized

@extent(1,3) LPTSTR s →

Valid

@valid @extent(2,3) LPTSTR s →
Validity: Related Work

• Validity is a lot like the Boogie methodology used in Spec#
  – Not as general since validity is just baked into macros
  – Many things are conditionally valid because of __success
    – Full conditional pre/post will allow more flexibility
• Even without it we can do some interesting with Objects
  – Treat them like structs!
  – Added in a few defaults
Structure Annotations

• Describes properties of buffers embedded in structs/classes

• Three scenarios supported
  – Outlined structure buffers
  – Structs with inline buffers
  – Header structs

• Structure descriptions interact with __in, __out, and __inout to determine pre/post rules for functions using structure buffers
struct buf {
    int n;
    __field_ecount(n)
    int *data;
};

struct ibuf {
    int n;
    __field_ecount(n)
    int data[1];
};

__struct_bcount(n * sizeof(int))
struct hbuf {
    int n;
    int data[1];
};
Zero Sized Buffers and NULL

```c
struct buf {
    int n;
    __field_ecount(n)
    int *data;
};
```

_opt versions available but generally not needed
class Stack {
public:
    Stack(int max);  // Stack(__out Stack *this,int max);
    int Pop();       // int Pop(__inout Stack *this);
    void Push(int v); // void Push(__inout Stack *this,int v);
    ~Stack();  // treated specially
private:
    int m_max;
    int m_top;
    __field_ecount_part(m_max,m_top)
    int *m_buf;
};
Conclusions

• Developers will accept the use of appropriate light weight specifications!
• But must understand the problem and tailor custom solutions
• Generic recipe:
  1) Write the problem down.
  2) Think real hard.
  3) Write the solution down.
  4) Repeat!
Questions?